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Request for grant of a patent

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1.	Your reference	UK9-98-093		
2.	Patent application number (The Patent Office will fill in this part)			
3.	Full name, address and postcode of the or of each applicant (underline all surnames)	INTERNATIONAL BUSINESS MACHINES CORPORATION Armonk New York 10504 United States of America		
	Patents ADP number (if you know it)	6537583 001		
	If the applicant is a corporate body, give the country/state of its incorporation	State of New York United States of America		
4.	Title of the invention	VIRTUAL MACHINE MEMORY MANAGEMENT		
5.	Name of your agent (if you have one)	P Waldner		
	"Address for Service" in the United Kingdom to which all correspondence should be sent (including the postcode)	IBM United Kingdom Limited Intellectual Property Department Hursley Park Winchester Hampshire SO21 2JN		
	Patents ADP number (if you know it)	6379853001		
6.	If you are declaring priority from one or more earlier patent applications, give the country and the date of filing of the or of each of these earlier applications and (if you know it) the or each application number	Country	Priority App No (if you know it)	Date of filing (day/month/year)
7.	If this application is divided or otherwise derived from an earlier UK application, give the number and the filing date of the earlier application	No of earlier application	Date of filing (day/month/year)	

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8. Is a statement of inventorship and of right to grant of a patent required in support of this request? (Answer 'Yes' if:  
a) any applicant named in part 3 is not an inventor, or  
b) there is an inventor who is not named as an applicant, or  
c) any named applicant is a corporate body.)

Yes

9. Enter the number of sheets for any of the following items you are filing with this form. Do not count copies of the same document

Continuation sheets of this form

Description

12

Claim(s)

3

Abstract

1

Drawing(s)

6

10. If you are also filing any of the following, state how many against each item.

Priority documents

-

Translations of priority documents

-

Statement of inventorship and right to grant of a patent (Patents Form 7/77)

3

Request for preliminary examination and search (Patents Form 9/77)

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Request for substantive examination (Patents Form 10/77)

Any other documents (please specify)

Fee sheet (FS1)

11. I/We request the grant of a patent on the basis of this application

P. Waldner  
Signature

21 December 1998  
Date

12. Name and daytime telephone number of person to contact in the United Kingdom

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Statement of inventorship and of  
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- |    |   |   |
|----|---|---|
| 1. | Your reference  | UK9-98-093                                  |
| 2. | Patent application number<br>(if you know it)   |   |
| 3. | Full name of the or of each applicant   | INTERNATIONAL BUSINESS MACHINES CORPORATION |
| 4. | Title of invention  | VIRTUAL MACHINE MEMORY MANAGEMENT           |
| 5. | State how the applicant(s) derived the right from the inventor(s) to be granted a patent  | By employment and by agreement              |
| 6. | How many, if any, additional Patents Forms 7/77 are attached to this form?  |   |
| 7. | I/We believe that the person(s) named over the page (and on any extra copies of this form) is/are the inventor(s) of the invention which the above patent application relates to. |   |

David Williams

Signature

21 December 1998  
Date

- |    |  |                           |
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| 8. | Name and daytime telephone number of person to contact in the United Kingdom | P WALDNER<br>01962 816057 |
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If there are more than three inventors, please write their names and addresses on the back of another Patents Form 7/77 and attach it to this form

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DUPLICATE

## VIRTUAL MACHINE MEMORY MANAGEMENT

FIELD OF INVENTION

5           This invention relates to memory management in a multithreaded runtime environment.

BACKGROUND OF INVENTION

10           The Java programming language has its origins in a project undertaken by Sun Microsystems to develop a robust programming environment that would meet the technical challenges of the consumer device software environment. The original consumer device projects were eventually abandoned but the Java programming language found itself being  
15           used on the World Wide Web to enable cross platform operation of programs downloaded from the internet. It is simple to use having similar features to C++ such as the basic object orientated technology but without some of the more complex features.

20           Typically, Java applications (source code) are compiled by the Javac compiler into Java byte code (intermediary code or pseudo object code) which can be loaded and executed by a Java Virtual Machine (JVM) (see Figure 1). The JVM provides an instruction set, memory management and program loading capability that is independent of the hardware  
25           platform on which it has been implemented. The Java application source code is compiled into architecture independent byte code and the byte code is interpreted by a JVM on the target platform. Java is designed to be portable and follows some defined portability standards, which intend the source code to be "write once, run anywhere". The Java byte code may  
30           be further compiled into machine code (object code) for the target platform at which point the architectural independent nature of Java is lost.

          The JVM is a software computing machine, effectively it simulates a hardware machine that processes Java byte code. The byte code is  
35           interpreted and processed by a JVM such as an Windows JVM running on a Intel personal computer platform. The JVM includes components for loading class files, interpreting the byte code, garbage collecting redundant objects, and for managing multiple processing threads. The JVM may also include a Just-In-Time compiler to transform some or all the  
40           byte code into native machine code.

          Multithreading is a feature built into the Java language to allow users to improve interactive performance by allowing operations to be performed while continuing to process user actions. Multithreading is  
45           similar to multitasking, but whereas multitasking allows many applications to run on the same system in several processes.

multithreading allows many routines (threads) in one application to potentially run in parallel within one process.

Garbage collection is the term used for describing how program objects are automatically discarded by the system after they are no longer useful.

For further information on garbage collection see Chapter 1 of 'Garbage Collection' by H Jones & R Lins, Wiley. Chapter 4 deals with Mark & Sweep techniques.

Many current implementations of Java use the classic mark-sweep-compact method of garbage collection as delivered in the base SUN JVM. References to the objects that are being processed at any instant by the system are stored in the registers, one or more thread stacks and some global variables. The totality of objects that are maybe needed by the system can be found by tracing through the objects directly referenced in the registers, stacks, and global variables and then tracing through these "root" objects for further references. The objects in use by a system thereby form a graph rooted at the root objects and any extraneous objects are not part of this graph. Once all the objects in the graph are found, the remaining objects in the heaps may be discarded (garbage collected).

The traditional mark and sweep garbage collection method is described below in terms of pseudo code with respect to a single heap:

- Stop all threads causing the active registers for each thread to be stored in its stack
- Trace all stacks for object references - the local roots
- Trace all global variables object references - the global roots
- Trace through root set for references until no new object references (the sum of the local and global roots is the root set).
- Delete all objects in the single heap that are not referenced

Note that in a particular Java virtual machine the term 'global variables' may include variables in classes, the global JNI table, objects awaiting finalisation and strings which have been 'interned'.

There are problems with this technique in a multi-threaded and long running environment. The first problem is that in order to garbage collect all the threads must be stopped in order to work out what

objects are unreachable (i.e., there are no pointers to them in the global or local variables and no pointers to them in other reachable objects). Various authors have attempted to solve this problem. One approach is an on-the-fly collector, which does not stop all threads, however, it cannot compact the reachable objects leading to fragmentation. Another approach are the generational scavenging schemes, which is to reduce the size of the set of traced objects by concentrating effort on the most recently allocated objects; however, these schemes must stop all of the threads. In an ideal world we would like to achieve a collector which works independently on all threads and compacts the local heap of the threads to maximise the free space available.

Another solution attempts to achieve this in a language (ML) other than Java by taking advantage of immutable objects which can be placed in thread-local heaps. An immutable object is non-modifiable and when such an object become reachable globally a copy of the object can be made in the global heap. Clearly this technique is only applicable to languages defining immutable objects.

Another approach moves an object into the global heap on first use. The difficulty here is that in order to move the object, references from elsewhere to it must be updated; in an environment where objects are referenced by handles this is made easier although there are still cases where objects cannot be moved. Unfortunately handles bring their own problems and the IBM ports of the JVM have removed handles to improve performance and remove the need to subdivide the heap into handles and object spaces.

#### SUMMARY OF INVENTION

According to one aspect of the invention there is provided a method of managing memory in a multi-threaded processing environment including respective local thread stacks and heaps and a global heap, said method comprising:

creating an object in a thread heap; and monitoring whether the object is reachable from anywhere other than the thread stack or registers.

Preferably the method further comprises: associating a local status with the object; and changing the status of the object to global under certain conditions.

More preferably the method further comprises deleting from the thread heap one or more local objects when they are not reachable from a local root.

Advantageously where reachability is determined by tracing from the local root.

5 More advantageously the status of an object in the thread heap is changed to global if the object is assigned to a static variable or if the object is assigned to a field in a global object.

10 Hence this proposal concentrates on the case of implementing a thread local heap which has not been done before. In essence the solution is to keep track of object references using write barriers placed in any operation which assigns references eg putfield, putstatic and astore.

15 The putfield operation for example is a java bytecode which causes the value on the top of the stack to be placed in the object referenced below it on the stack. The object field to be updated is defined by the constant pool reference (2 bytes) which follows the bytecode. See p.325 'Java Virtual Machine Specification', Lindholm & Yellin 1997.

20 A flag is associated with each object in the heap. In practice the flag is easily provided by ensuring that all objects on the heap are allocated on an 8 byte boundary ensuring that the low three bits are 0. Two of these bits are currently used to denote pinned or free objects leaving a spare flag which we use to denote 'global'. When an object is  
25 allocated it is generally placed in a piece of storage associated with the allocating thread; the thread local heap. The global flag is unset at this point. From now on the write barriers implement the following rules:

30 1) if an object reference is assigned to a static (global) variable the object becomes global and all objects reachable from it also become global; and

35 2) if a reference to a non-global object is assigned to another object which is global then the referenced object and all objects that are reachable from it also becomes global.

40 Thus at any point the thread local heap will contain objects which have the global flag unset and some which have the flag set. Any objects with the flag unset can be garbage collected by the thread itself. The objects marked global must at this stage be treated as reachable, pinned objects; they must not be moved or deleted. Any garbage collection policy can be adopted including compaction for the objects in the local heap, the only constraint is that we must of course compact into the  
45 potentially fragmented spaces left between the global objects. Note that

the write barriers described above are extremely simple and can be implemented to have a very small effect on performance.

At some point the global objects will be moved from the local heaps into the global heap and this must be done using the current collection approach of stopping all threads. The belief is of course that the frequency of such events can be dramatically reduced and the work to be done also reduced. The global objects need to be removed and put in their own heap, once that is done the individual threads can be left to continue any tidying they need to do on their own heaps. The great improvement is that the moving of global objects is the only part which must be done whilst all program threads are stopped. Tidying of the local heaps can be then done in parallel. The current approach forces the compaction of the entire heap (moving objects) to be done by one thread and all others must wait for completion.

#### BRIEF DESCRIPTION OF DRAWINGS

In order to promote a fuller understanding of this and other aspects of the present invention, an embodiment will now be described, by way of example only, with reference to the accompanying drawings in which:

Figure 1 is a schematic representation of a platform supporting the embodiment of the present invention;

Figure 2 is a schematic representation of a Java Virtual Machine embodying the invention;

Figure 3 shows the structure of objects in a thread heap;

Figure 4 is a flow diagram depicting the memory management process of the invention;

Figure 5 is a flow diagram of the process of the write barrier; and

Figure 6 is a schematic representation of objects in a thread heap and the global heap of the JVM.

#### DETAILED DESCRIPTION OF PREFERRED EMBODIMENT

Referring to Figure 1 there is shown a computer platform 10 such as pentium processor based PC typically comprising motherboard with processor, 32Mbytes RAM memory, and interface circuitry; 4G Byte hard drive; 20x CD ROM; SVGA monitor; keyboard and mouse. On power on the platform loads into the memory an operating system 12 such as Windows NT

v4 and subsequently a Java Virtual Machine (JVM) based on the Java Development Kit (JDK) v1.1 or v1.2. The Java virtual machine is capable of loading a Java application 16 from disk into memory so that the application may be executed on the platform. The Java application 16 is object orientated program code and creates Java language objects 18A, 18B, 18C, 18D. An example of how the objects are built by the JVM is described with reference to Figure 3.

The JVM 14 comprises code components for the multitude of tasks it has to perform (see Figure 2). The components of relevance to the present embodiment comprise memory manager 20, Java interpreter 22 and just-in-time compiler 24. In operation the JVM creates memory space, the runtime data area 26, for storing the objects 18 and data for the application 16.

The runtime data area 26 comprises thread memory space 28A, 28B, 28n, global heap 34, class area 36 and compiled method area 38. The thread memory space 28 ideally stores local thread objects, local thread execution data and also objects marked global. Local objects are objects which are used by one thread alone and are not common between. Thread objects are stored in thread heaps 32A, 32B, 32n and thread execution data including object references and register data are stored in the thread stacks 30A, 30B, 30n. The global heap 34 stores objects which are common to more than one thread, that is global objects. A local object which is loaded into a thread heap 32A for processing by thread 28A may become global if it is later used by another thread, say thread 28B, and may remain in the thread heap 28A or be moved to the global heap 34 by 'stop everything'. Class area 36 stores the classes as they are loaded into the JVM, classes are defined as global because they are common to all the threads hence any field of a class is globally reachable. An object, which is an instance of a class, is placed in a thread heap and is initially local because it may only be used by that thread. An exception is instance objects of class Thread or Runnable which are global at creation because they can be reached by other threads. Whereas an object referenced by a variable in a class is global because the object is reachable from the class which is reachable by all threads. The compiled method area 38 is used by the JIT compiler 24 to store native executable code compiled from the Java byte code.

The memory manager 20 includes an object allocator 21A and a garbage collector 21B which between them control the creation and deletion of objects within the thread heaps 32A, 32B, 32n and the global heap 34.

The Java interpreter 22 includes functionality to process the multitude of Java byte codes which make up the Java byte code function



set. A write operation component 23A processes java write operations such as 'putfield', 'putstatic', 'aastore' using the operating system and microprocessor of the platform. Putfield and aastore are write operations which among other things assign a source object reference to a field in a target object. This means that the target object will contain a reference (or pointer) to the source object; the target object is the object that will be updated by the putfield operation. Putstatic assigns an object reference to a static variable in a class. A write barrier component 23B checks the status of the target object and assigns a local or global status to the source object depending on certain rules:

1) if an object reference (of a source object) is assigned to a static variable of a class the source object becomes global and all objects reachable from it also become global; and

2) if a non-global source object is assigned to a reference in a target object which is global then the referenced source object and all objects that are reachable from it become global.

The Java write operation 'putfield' may set a field within a target object to point at a source object using the operating system commands or platform operations. A 'putfield' write operation 23A together with the write barrier 23B is described below in pseudo platform code with comments.

If field to be updated in target object is of 'reference' type.

if (target is global) {Check to see to the global flag is set for the target

if (source is not global) {If so check to see if the source has already been set global

set source to global {If not already set then set source to global

set objects reachable by source to global

{and set all objects reachable,  
(i.e., directly referenced by source  
object or referenced from another  
object reachable from the source  
object by the source object to global

set slot (target, field number, source)

{sets a field variable within a target object to point at a source object.

5           The last instruction code is the write operation 23A whereas the instructions before that constitute the write barrier 23B. Similar pseudo code may be provided for the putstatic and aastore java commands, and other JVM actions, which store references in objects.

10           The JIT compiler 24 is similar to the Java interpreter 22 in that it contains a write operation component 25A and a write barrier component 25A. However instead of interpreting java byte code in the methods of objects and processing the operations in real time, native code similar to the pseudo code above is created and stored in the compiled method  
15           area 38 for execution directly by the platform and operating system whenever that particular method is invoked.

          Objects 18 are stored in the thread heaps 32 using multiples of 8 bytes (see Figure 3). Each object has a length word attached which  
20           identifies the length of the object in bytes and allows objects in the heaps to be scanned sequentially from the start of the heap. The length word (4 bytes long in the embodiment) has three spare bits available because of the 8 byte alignment boundary and one of these bits is used as a flag for the object to store the local or global status. In the  
25           embodiment shown in Figure 3 each of the objects is 40 bytes long which is 5\*8 bytes or '00101000' bytes in binary; the last three '000' being spare. This is indicated at the high end of the word in each case. The low end of the word is spare and thus the first bit of the length word is used to indicate the local/global status, in this case object 18A and 18B  
30           are set to '1' to indicate that they are global and 18C is set to '0' to indicate that it is local.

          The process of memory management is described with reference to Figure 4. Step 4.1 an object is created from a class stored in the class  
35           area 36. Step 4.2 check to see if the class is a 'global' class, i.e., a class all of whose instances are global or one whose instances we expect to quickly become global, whereby the object is assigned global and/or placed in the global heap (Step 4.7). Step 4.3 the size of the object is calculated to see whether it will fit in the thread heap. Step 4.4 if  
40           there is not enough space then garbage collection is performed to free up memory. Step 4.5 place the object in the thread heap memory along with the object length in multiples of eight. Step 4.6 Use a spare bit in the length attribute as a flag and set as local ('0'). The process ends at  
45           step 4.8.

The write operation and write barrier process are described with reference to Figure 5. Step 5.1 a write operation is called and intercepted by the write barrier which is integrated with the write operation code. Step 5.2 Check whether the object is being assigned to a static variable of a class and if so set the status of the object as global (step 5.4) and set all objects reachable by the object as global (step 5.5) before moving to step 5.6, if not so then do step 5.3. Step 5.3 Check whether the write operation assigns the object to another global object, if so set the status of the object as global (step 5.4) and set all objects reachable by the object as global (step 5.5). Once complete continue with the write operation.

The garbage collector 21A for single thread garbage collection is described in terms of pseudo code with comments.

15	Trace stack in thread	{Identify objects in use by the thread - these are held in the thread stack
20	Trace reference objects	{Identify objects referenced directly and indirectly by the stack objects.
25	Delete unused objects	{those objects in the heap which are not identified in trace or are not identified as global
30	Check heap size	{If not enough memory then increase size of the local heap. If not desirable or possible then do 'full' garbage collection

Full garbage collection is traditional Mark/Sweep performed with all other threads stopped. During full garbage collection global objects are moved from local to the global heap if possible. Object movement is not possible if conservative tracing is performed and the object is referenced directly from the stack because of the uncertainty of the tracing.

The Java language pseudo code program below (together with comments) creates objects representing a bike with two wheels and a bell in a local thread which is put into a global shed. The bike and wheels are make global whereas the bell remains local. If a further new bike object were to be created and the thread heap had no more memory then the garbage collector would mark and sweep the bell object so long as it was not pointed to by the stack.

```
// Demonstrate the Thread local heap in action.
```

```
// For simplicity all niceties such as constructors have been  
omitted and the example
```

5

```
// made deliberately unreal to illustrate the basic point. Objects  
generally start life
```

```
// local and become global by assignment.
```

10

```
// This is a fake class just to illustrate a static variable in  
operation.
```

15

```
class Shed { // This is a fake class just to illustrate a static  
variable in operation.
```

```
static Bike;
```

20

```
// The main method will run the sample code to illustrate creation  
of local
```

```
// objects which later will become global.
```

25

```
public static void main(String args[]) {
```

```
Bike b = new Bike(); // build a new Bike object
```

```
// The Bike class is loaded and initialised  
and an instance of Bike is created in the local heap. It is marked  
'local'.
```

30

```
Bell c = new Bell(); // build a new Bell object. The Bell class  
is loaded and initialised and an instance of Bell is created in the local  
heap. It is marked 'local'.
```

35

```
Wheel w = new Wheel(); // build a new wheel object. The Wheel  
class is loaded and initialised and an instance of Wheel is created in  
the local heap. It is marked 'local'.
```

40

```
b.wheel1 = w; // Assign the first wheel of the bike Java  
bytecode putfield, write barrier detects that target (b) is local and  
hence w remains local.
```

45

```
Shed.bike = b; // Assign b to the static field in bike Shed.  
Java bytecode putstatic is used and thus b is marked global. Since w is  
reachable from b then w is also marked global.
```

w = new Wheel(); // build another new wheel object. Another instance of Wheel is created in the local heap. It is marked 'local'.

5           b.wheel2 = w;           // Assign the second wheel of the bike  
                                  // Java bytecode putfield, write barrier detects  
that target (b) is global and hence w is marked global also.

10           // Note that at this point we have built a bike, two wheels and a  
bell. The bell is the only object which is still local since the bike has  
been assigned to the static variable 'bike'.

```

    }
}
class Bike {
Wheel wheel1;
15  Wheel wheel2;
Bell bell1;
String name;
}
class Wheel {
20  int diameter;
}
class Bell {
int volume;
}

```

25           In summary there is described a method and system for memory  
management in a virtual machine or operating system and in particular an  
object creation and garbage collection method and system. There is  
described a method and system of managing memory in a multi-threaded  
30           processing environment such as a java virtual machine comprising:  
creating an object in a thread heap; associating a status with the object  
and setting the status as local; using write barriers to change the  
status to global if the object is assigned to a static variable or if the  
object is assigned to a field in a global object; changing the status of  
35           all objects referenced by that object to global; and performing garbage  
collection by deleting from the thread heap, when memory space in the  
thread heap is required, one or more local objects which can not be  
traced to the thread stack.

40           The platform in the embodiment does not need to be pentium based  
nor is restricted to the hardware stated. Any platform which is capable  
of supporting JDK v1.1 or v1.2 would be capable of supporting the virtual  
machine of embodiment. Although the embodiment is described with  
reference to a Java Virtual Machine it is not necessarily so restricted  
45           to a virtual machine but may be used in any environment where storage is  
garbage collected to free up memory in that environment. For example,

such an environment could be an multithreaded Lisp runtime or the runtime for some other multithreaded garbage collected language.

5           Although we have described in the preferred embodiment the local/global status as a field in the object it may be implemented in other ways for instance as a separate table.

10           Java and Java-based marks are trademarks or registered trademarks of Sun Microsystems, Inc.

15           Microsoft, Windows, and Windows NT are trademarks of Microsoft Corporation.

20           Now that the invention has been described by way of a preferred embodiment, various modifications and improvements will occur to those person skilled in the art. Therefore it should be understood that the preferred embodiment has been provided as an example and not as a limitation.

## CLAIMS

1. A method of managing memory in a multi-threaded processing environment including respective local thread stacks and heaps and a global heap, said method comprising:

creating an object in a thread heap; and

monitoring whether the object is reachable from anywhere other than the thread stack.

2. A method as claimed in claim 1 further comprising:

associating a local status with the object;

changing the status of the object to global under certain conditions.

3. A method as claimed in claim 2 further comprising deleting from the thread heap one or more local objects when they are not reachable from a local root.

4. A method as claimed in claim 3 where reachability is determined by tracing from the local root.

5. A method as claimed in claim 2,3 or 4 wherein the status of an object in the thread heap is changed to global if the object is assigned to a static variable or if the object is assigned to a field in a global object.

6. A method as claimed in claim 1, 2 or 3 further comprising intercepting assignment operations to an object in a thread heap to access whether the object status should be changed.

7. A method as claimed in any of the preceding claims wherein the multithreaded processing environment is a virtual machine.

8. A method as claimed in claim 7 wherein the virtual machine comprises an interpreter and the write operation code in the interpreter is modified to perform the checking of assignment of the object.

9. A method as claimed in claim 7 or 8 wherein the virtual machine comprises a just in time compiler wherein native compiled write operation code includes native code to perform the checking of assignment of the object.

10. A method as claimed in any one of the preceding claims further comprising using spare capacity in the object header for the flag.

5 11. A method as claimed in claim 10 further comprising using multiples of 2 or more bytes in a thread heap to store the objects whereby there is at least one spare bit in the object length variable and using the at least one spare bit as the flag.

10 12. A method as claimed in any of the preceding claims further comprising moving objects whose status is global from the thread heap to a global heap.

15 13. A method as claimed in any of the preceding claims further comprising compacting the reachable local objects in a thread heap.

14. A method as claimed in claim 1 wherein certain objects are associated with a global status on creation.

20 15. A method as claimed in claim 14 where said certain objects include class objects.

25 16. A method as claimed in claim 14 further comprising the step of analysing whether an object is likely to be made global and associating such an object with a global status on creation.

17. A method as claimed in claim 14, 15 or 16 further comprising allocating objects assigned as global on creation to the global heap.

30 18. A system for managing memory in a multi-threaded processing environment comprising:

respective local thread stacks and heaps;

35 a global heap;

means for creating an object in a thread heap; and

40 means for monitoring whether the object is reachable from outside the thread heap.

19. A system as claimed in claim 18 further comprising means for associating a local status with the object and means for changing the status of the object to global under certain conditions.



20. A system as claimed in claim 19 further comprising means for deleting from the thread heap one or more local objects when they are not reachable from a local root.

5 21. A system as claimed in claim 20 further comprising:

means for changing the status of an object in the thread heap to global if the object is assigned to a static variable or if the object is assigned to a field in a global object.

10 22. A computer program product stored on a computer readable storage medium for, when executed on a computer, performing a method of managing memory in a multi-threaded processing environment including respective local thread stacks and heaps and a global heap, said method comprising:

15 creating an object in a thread heap; and

monitoring whether the object is reachable from outside the thread heap.

20 23. A method as claimed in claim 22 further comprising:

associating a local status with the object;

25 changing the status of the object to global under certain conditions.

24. A method as claimed in claim 23 further comprising deleting from the thread heap one or more local objects when they are not reachable from a local root.

30 25. A method as claimed in claim 24 where reachability is determined by tracing from the local root.

35 26. A method as claimed in claim 23, 24 or 25 wherein the status of an object in the thread heap is changed to global if the object is assigned to a static variable or if the object is assigned to a field in a global object.

## VIRTUAL MACHINE MEMORY MANAGEMENT

## ABSTRACT

5           This invention relates to memory management in a virtual machine or  
operating system and in particular to object creation and garbage  
collection. There is described a method and system of managing memory in  
a multi-threaded processing environment such as a java virtual machine  
comprising: creating an object in a thread heap; associating a status  
10 with the object and setting the status as local; using write barriers to  
change the status to global if the object is assigned to a static  
variable or if the object is assigned to a field in a global object;  
changing the status of all objects referenced by that object to global;  
and performing garbage collection by deleting from the thread heap, when  
15 memory space in the thread heap is required, one or more local objects  
which can not be traced to the thread stack.



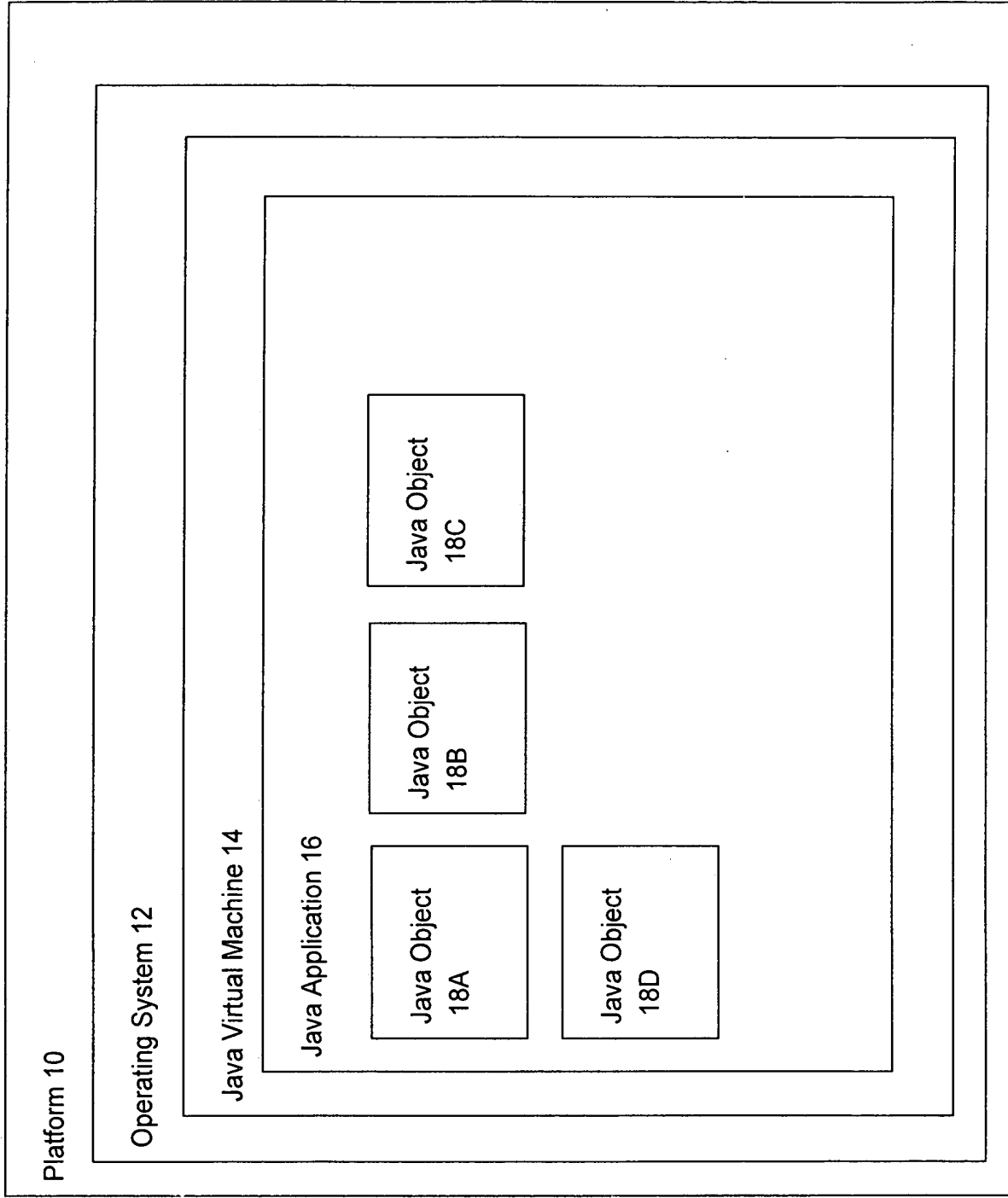


Figure 1

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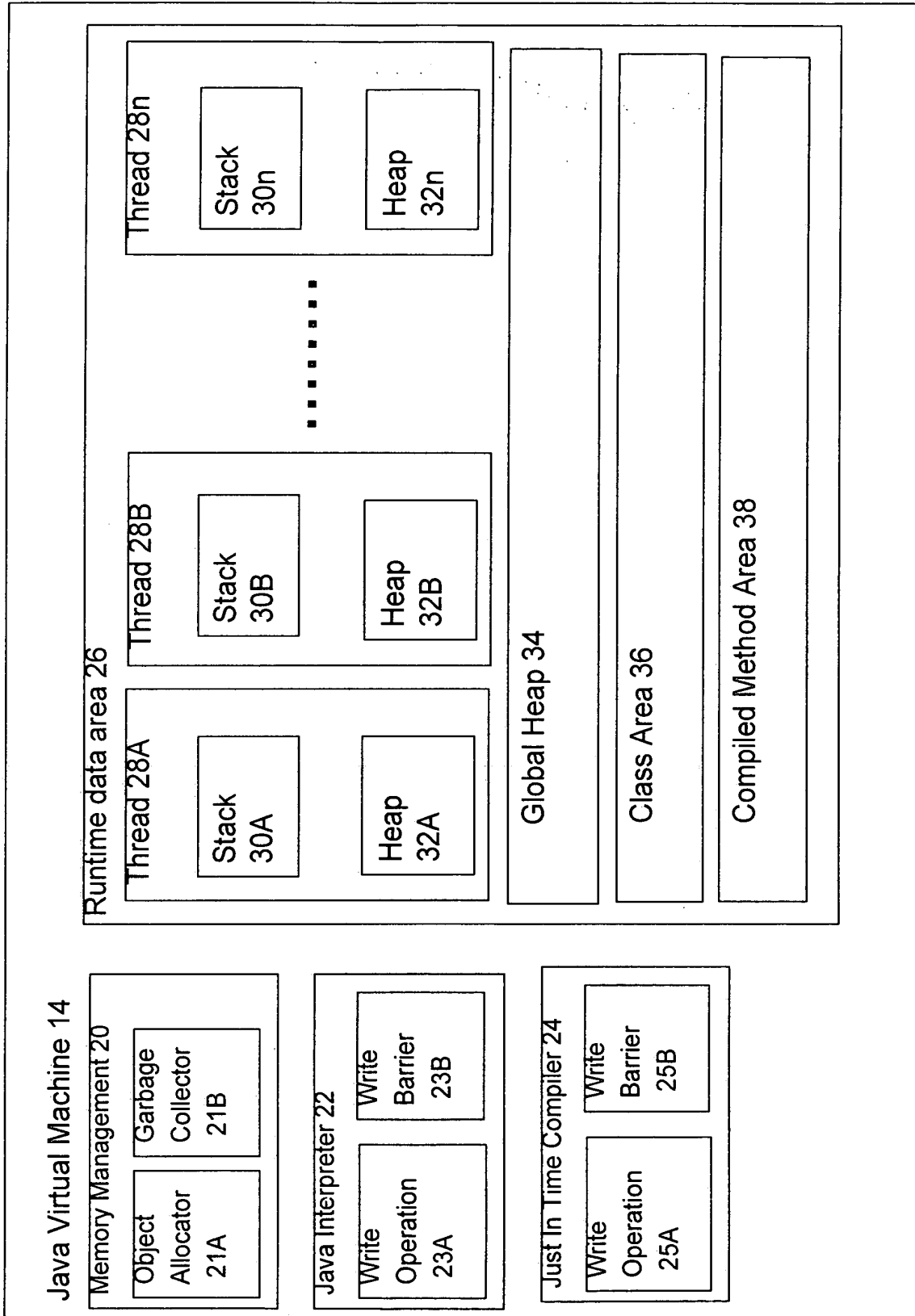


Figure 2

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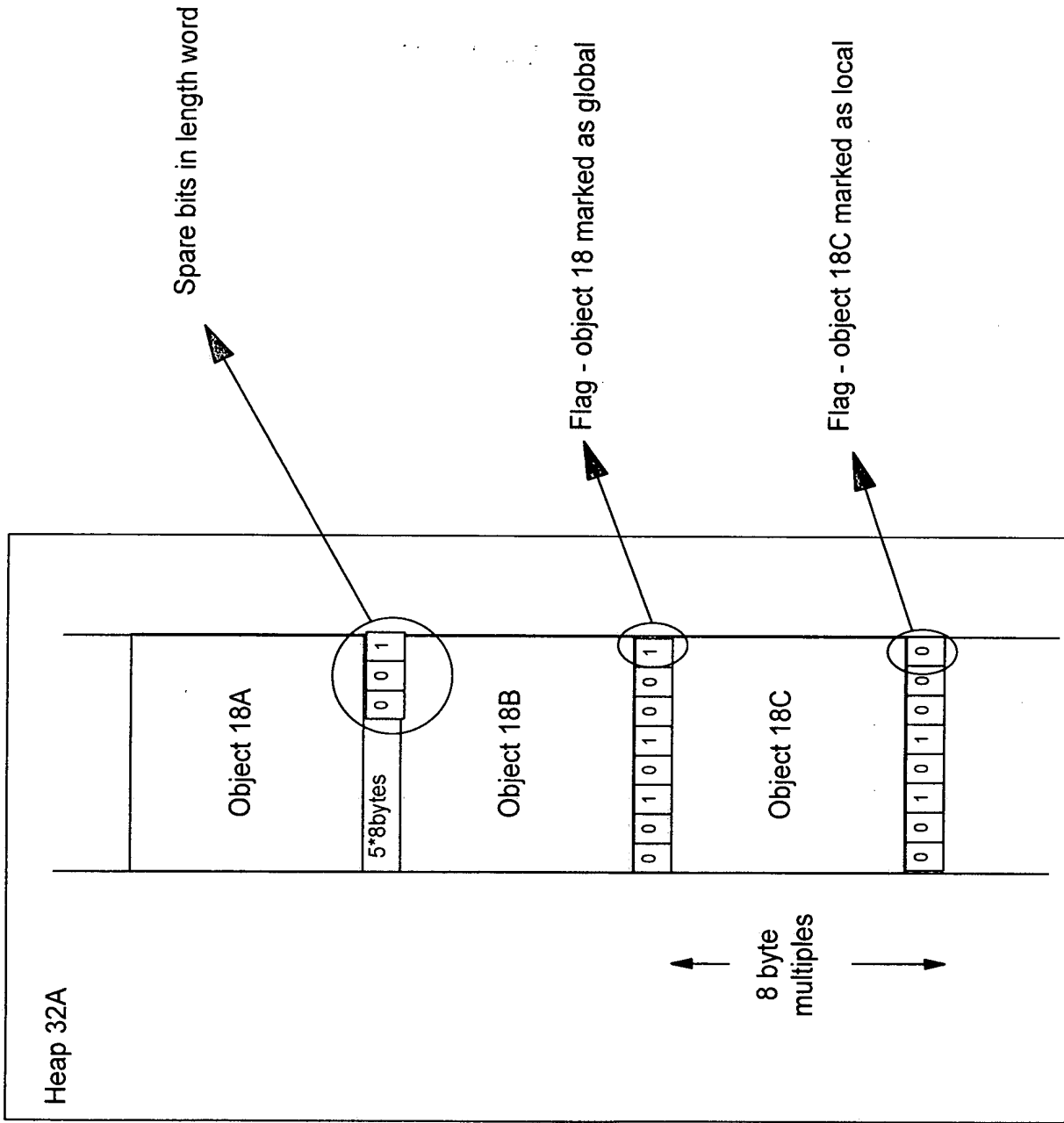


Figure 3

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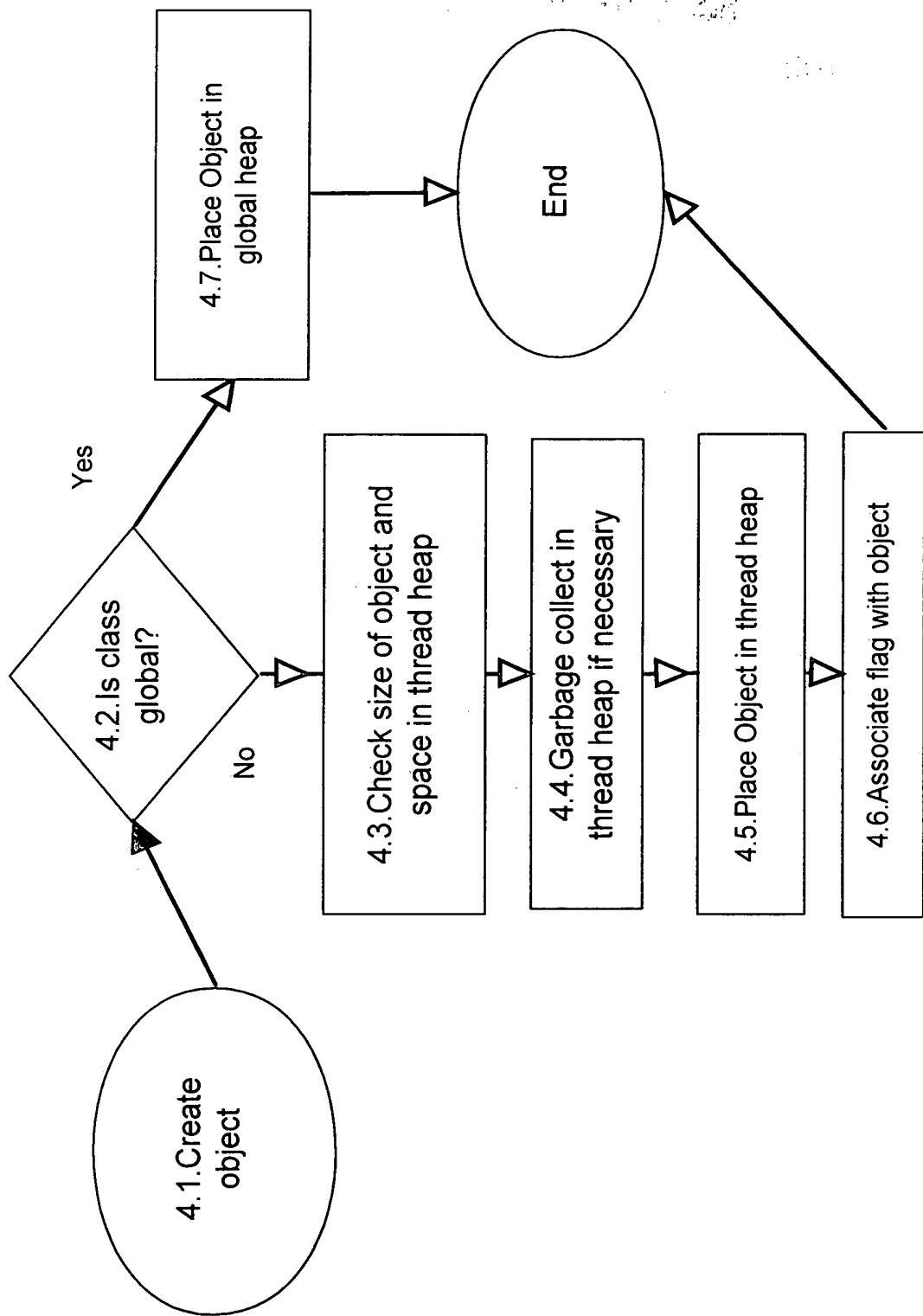


Figure 4

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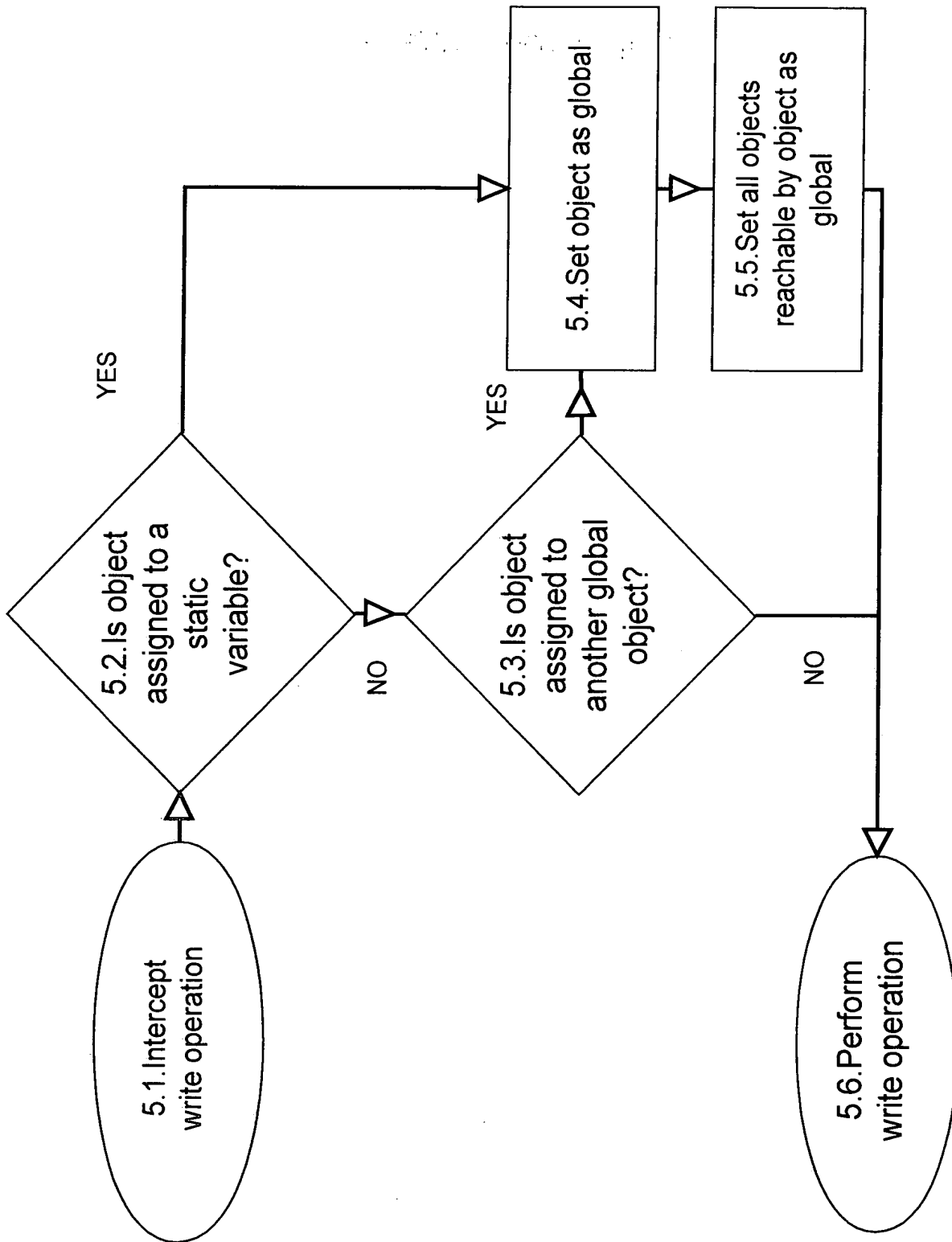


Figure 5

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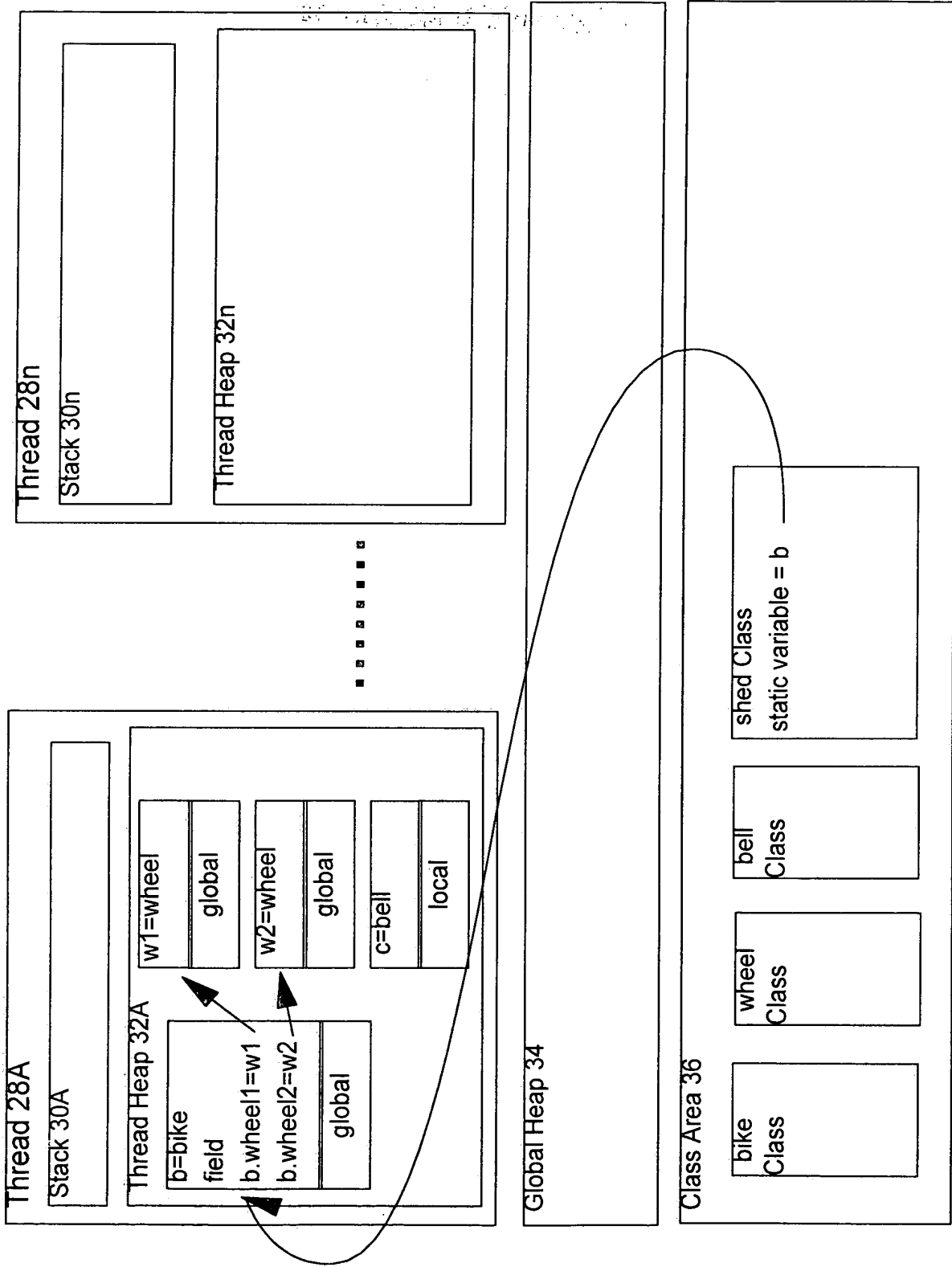


Figure 6

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